pidfds: Process file descriptors on Linux

Tuebix, 2019 Tübingen, Germany

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christian@brauner.io christian.brauner@ubuntu.com <u>@brau_ner</u> <u>https://brauner.io/</u> <u>https://people.kernel.org/brauner</u> Processes, Threads, PIDs, TIDs, fds, and other assorted nonsense

A pidfd is a file descriptor referring to a process or a thread*.

* At some point in the future

Processes

- dumb definition would be "executing instance of a program"
- a binary image
- memory (stack, heap)
- resources (file descriptors etc.)
- attributes (uid, gid, capabilities, etc.)
- fork() + exec() to spawn new process load binary image

```
struct linux_binprm
exec_binrprm()
-> search binary handler(<aout, elf, elf fdpic, em86, flat, script>)
```

Processes

What does fork() do?

- **duplicates/copies** the calling process
- duplicated/copied process -> parent
- duplicate/copy -> child
- typical code pattern (abbreviated at the end):

```
int ret;
pid_t pid, wpid;

pid = fork();
if (pid < 0)
    _exit(EXIT_FAILURE);

/* child */
if (pid == 0) {
    printf("I am the child");
    _exit(EXIT_SUCCESS);
}

/* parent */
wpid= waitpid(pid, NULL, 0);</pre>
```

Threads

- lightweight process
- processes exclusively own their resources (for the most part)
- threads share their resources (for the most part)
- clone() to create a new thread by specifying a bunch of flags; at least:
 CLONE_THREAD, CLONE_FILES, CLONE_FS, CLONE_SYSVSEM, CLONE_SIGHAND,
 CLONE_SETTLS, CLONE_VM
- userspace library (part of libc) to manage threads: pthreads

Processes and Threads

- vague terminology
- kernel usually uses "task" which is best translated to "thread"
- single thread ("single-threaded") -> process
- multi-threaded -> multiple threads with one thread-group leader

PIDs, TIDs, and for the sake of confusion TGIDs

- PID -> process identifier
- TID -> thread identifier
- TGID -> thread-group identifier
- single threaded processes PID == TID == TGID
- multi-threaded process TGID == PID != TID for non-thread-group leaders
- thread-group leader: first process that created a new thread
- thread-group leader stays zombie until all other threads in its thread-group have exited (that has various interesting consequences... for some definition of "interesting")

Signals

- Signals are one way for userspace to interact with processes and threads
- Can also be used for IPC, i.e. can be used to communicate between processes (or threads)
- Signals can be targeted at a whole thread-group or at a single thread.

Killing the wrong target

You really wanted to kill that pesky little resource hungry bastard but instead you killed a really essential little sucker.

PID allocation

- PIDs are allocated cyclically alloc_pid()
 -> idr_alloc_cyclic()
 The standard PID limit is 32768 but can be up to 2^22

PID recycling

- PID recycling is built into the allocation algorithm
- on high pressure systems it's rather easy to operate on the wrong process...
- possible mitigations:
 - make PID allocation algorithm random
 - bump limit to 2^22
 - introduce UUIDs
- "but either way, I don't like correctness guarantees based on timing, especially if they might affect security"

pidfds

- file descriptors:
 - handle on a file {file, directory, device, etc...}
 - private, stable reference
- skipping over a few details, an fd references a struct file in the kernel
 - multiple fds can refer to the same struct file
 - multiple struct file can refer to the same inode

```
int fd1 = open(something);
int fd2 = dup(fd1);
int fd3 = open(something);
```

pidfds

- every struct file comes with a member called f_op which refers to a struct file_operations
 - f_op essentially defines what "type" a file is standard pattern in kernel comparing file operations to determine type of file, e.g. file->f op == &pidfd fops
 - struct file also comes with another member called private_data
 - used to stash file specific data, e.g. pidfds stash away a struct pid (the kernel's abstraction on top of tasks
- struct pid is the kernel-internal stable handle on a process

pidfds: A New Hope^wAPI

- building a new process management API
- at some point you can manage processes completely without using PIDs
- Linux 5.1 (released)
 - pidfd send signal()
- Linux 5.2 (to be released tomorrow or next Sunday)
 - CLONE PIDFD
- Linux 5.3 *(upcoming merge-window)*
 - pidfd_open()
 - polling support
- Linux 5.4
 - WPIDFD

Can we do more? Maybe...;)

pidfds: User? User!

- systemd
- Android's LMKD
 Google is already backporting the complete pidfd work
 to all of their LTS kernels:
 - 4.9: https://android-review.googlesource.com/q/topic:%22pidfd+polling+support+4.9+backport%22
 - 4.14: https://android-review.googlesource.com/q/topic:%22pidfd+polling+support+4.14+backport%22
 - 4.19: https://android-review.googlesource.com/q/topic:%22pidfd+polling+support+4.19+backport%22
- LXC and LXD

pidfds: A Linux invention?

- No!
- Solaris and the BSDs have them too

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